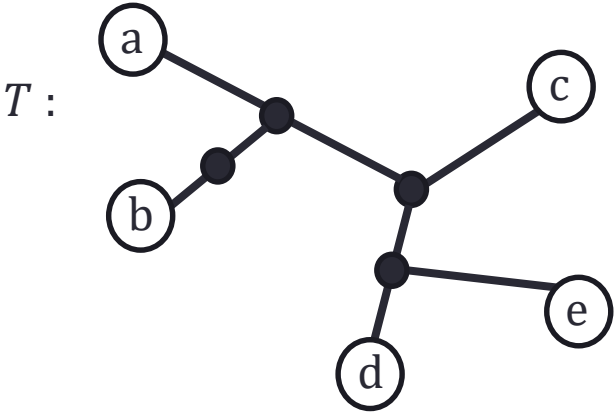


WILL WE EVER FIND A FORBIDDEN SUBGRAPH CHARACTERIZATION OF LEAF POWERS?

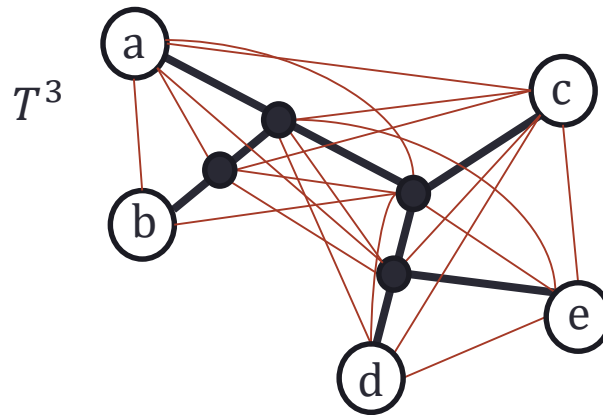
Manuel Lafond, Université de Sherbrooke, Canada



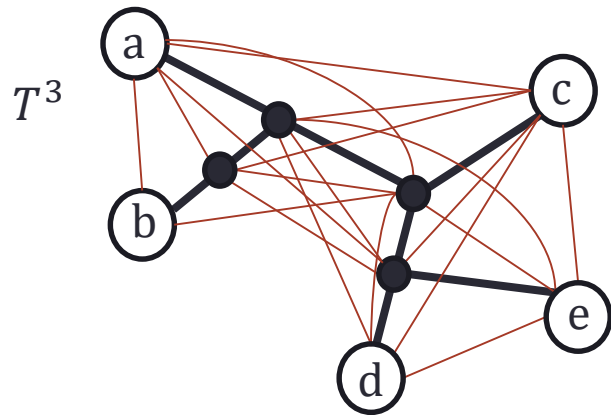
Tree T , leafset denoted $L(T)$, leaves have distinct labels



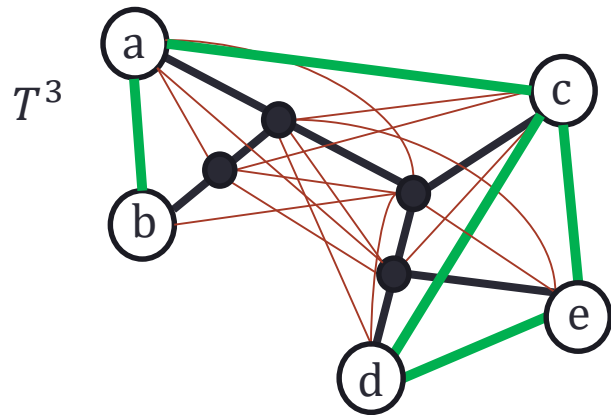
- T^k : k -th power of T
 - $V(T^k) = V(T)$
 - $uv \in E(T^k) \Leftrightarrow \text{dist}_T(u, v) \leq k$



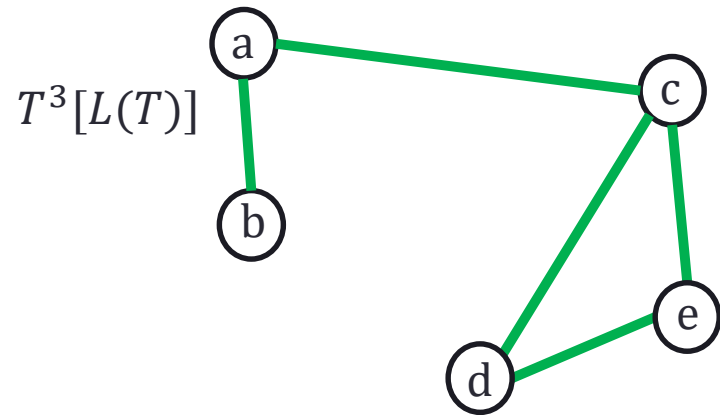
- The **k -th leaf power** of T is $T^k[L(T)]$, the subgraph induced by the leaves.



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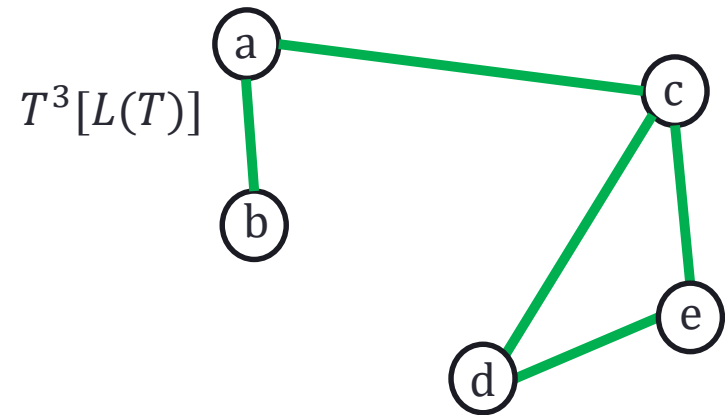
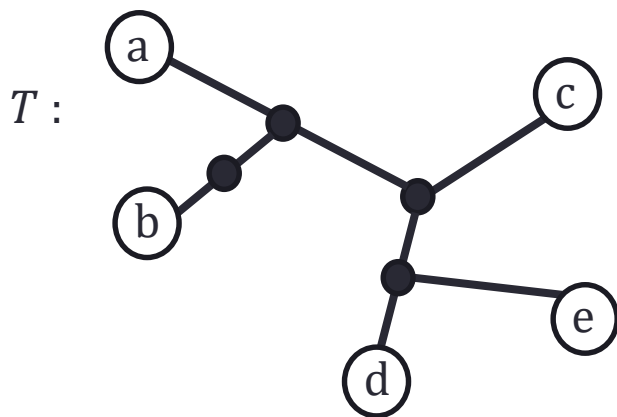


- The ***k*-th leaf power** of T is $T^k[L(T)]$, the subgraph induced by the leaves.



Definition

A graph G is a **k -leaf power** if there exists a tree T such that $G = T^k[L(T)]$



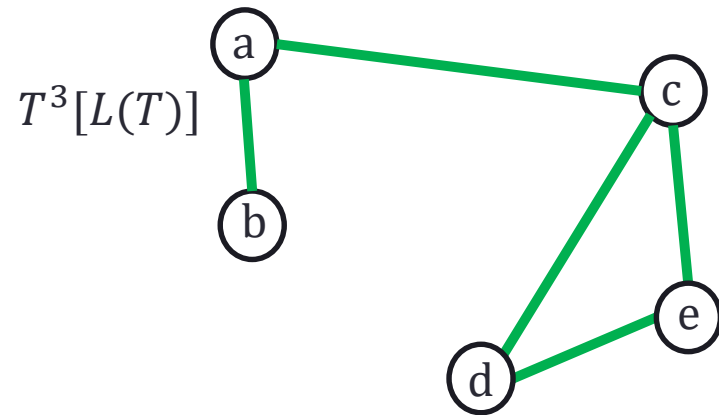
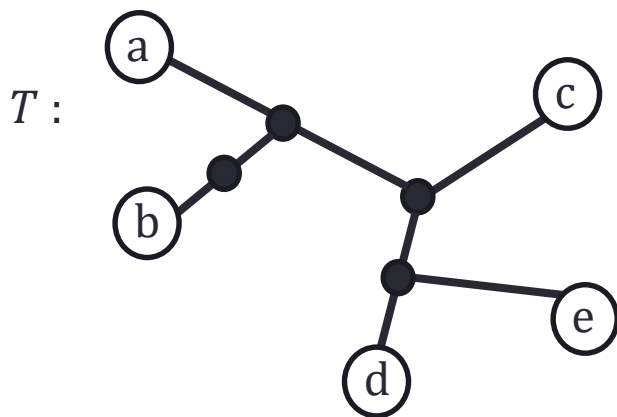
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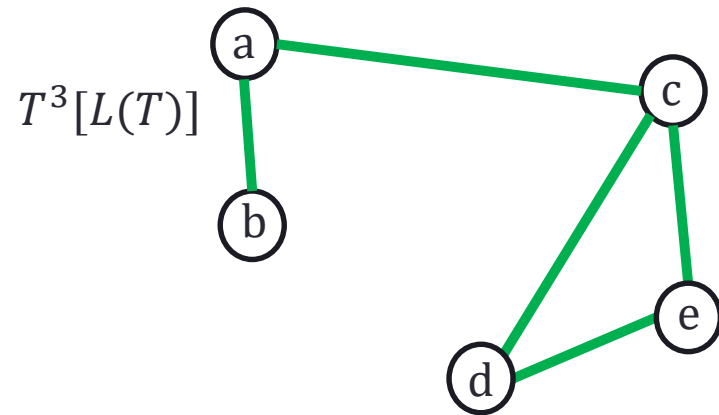
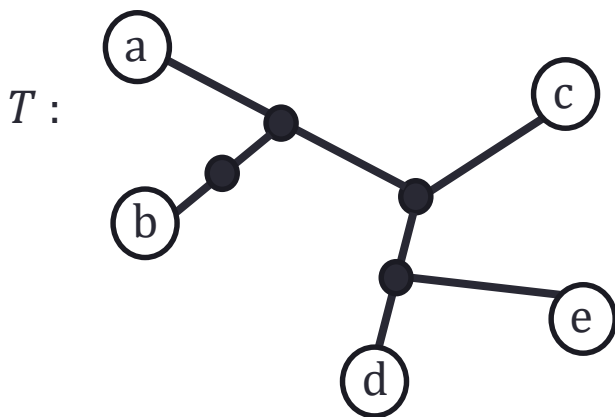
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Definition

A graph G is a **leaf power** if there exists k such that G is a k -leaf power .

Graph theoretical problems [Nishimura, Ragde, Thilikos, 2002]

- Provide a forbidden induced subgraph characterization of leaf powers.
- Provide a forbidden induced subgraph characterization of k -leaf powers, for every k .

Algorithmic problems

- Is deciding whether a graph is a leaf power in P?
- For fixed k , is deciding whether a graph is a k -leaf power in P?

In this talk

Graph theoretical problems

- Provide a forbidden induced subgraph characterization of leaf powers.

Conjecture (Nevries & Rosenke, 2015)

There are only a finite number of minimal strongly chordal graphs that are not leaf powers.

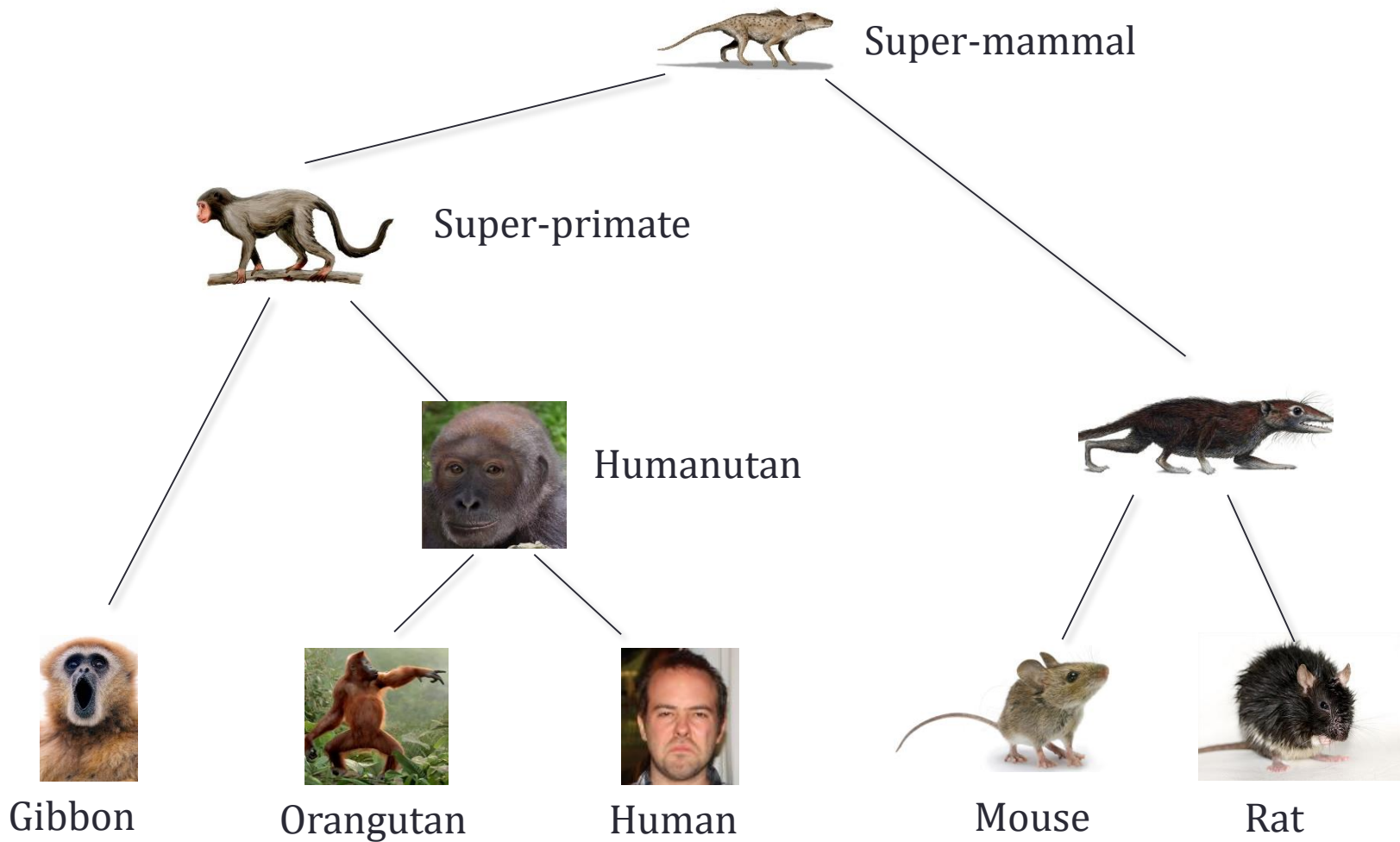
(implication: easy poly-time algorithm)

(L, 2018)

There is an infinity of minimal strongly chordal graphs that are not leaf powers.

(implication: still no easy poly-time algorithm)

Motivation



Motivation



Gibbon



Orangutan



Human



Mouse



Rat

Motivation

- Distance matrix.
- Noisy.
- Basic info: "near" or "far"
 - w.r.t some threshold

	G	O	H	M	R
G	0	0.33	0.25	0.52	0.61
O		0	0.31	0.61	0.36
H			0	0.55	0.56
M				0	0.2
R					0



Gibbon



Orangutan



Human



Mouse



Rat

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M				0	0.2
R					0

threshold = 0.5



Gibbon



Orangutan



Human



Mouse



Rat

Motivation

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- Basic info: "near" or "far"
 - w.r.t some threshold

	G	O	H	M	R
G	0	1	1	0	0
O		0	1	0	1
H			0	0	0
M				0	1
R					0



Gibbon



Orangutan



Human



Mouse

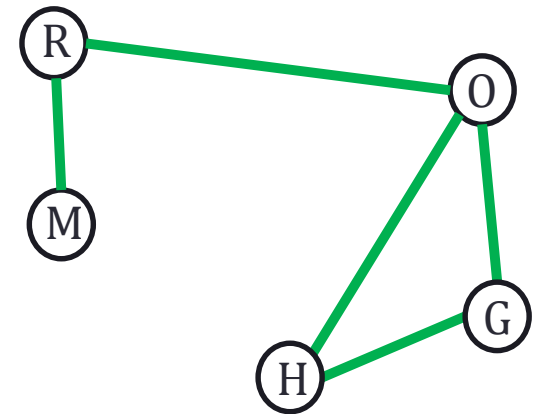


Rat

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H			0	0	0
M				0	1
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Gibbon



Orangutan



Human



Mouse



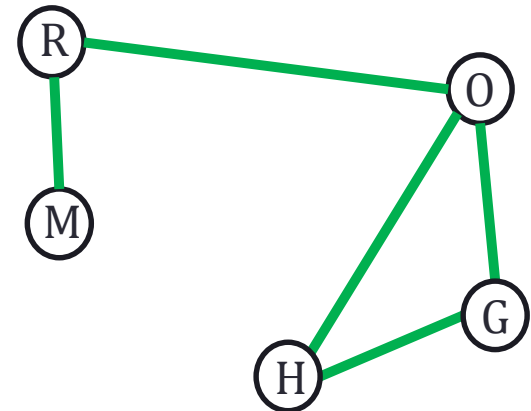
Rat

Motivation

- This graph "makes sense" if we can find a tree T in which neighbors are "close", non-neighbors are "far".
- Does there exist k and a tree T with leafset $V(G)$ satisfying

$$uv \in E(G) \Leftrightarrow \text{dist}_T(u, v) \leq k$$

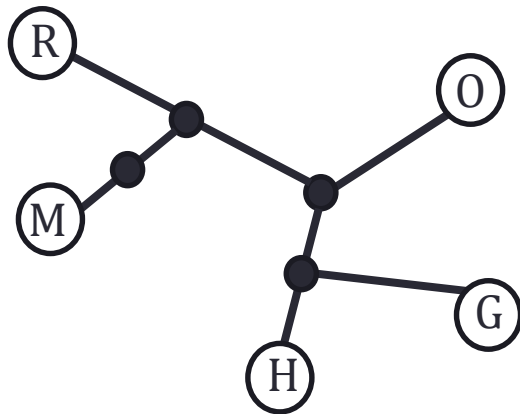
	G	O	H	M	R
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M				0	1
R					0



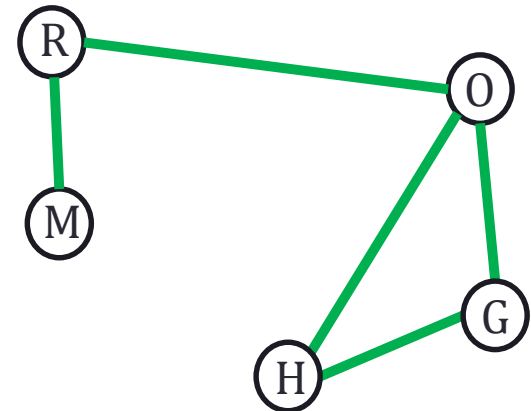
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H			0	0	0
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R					0



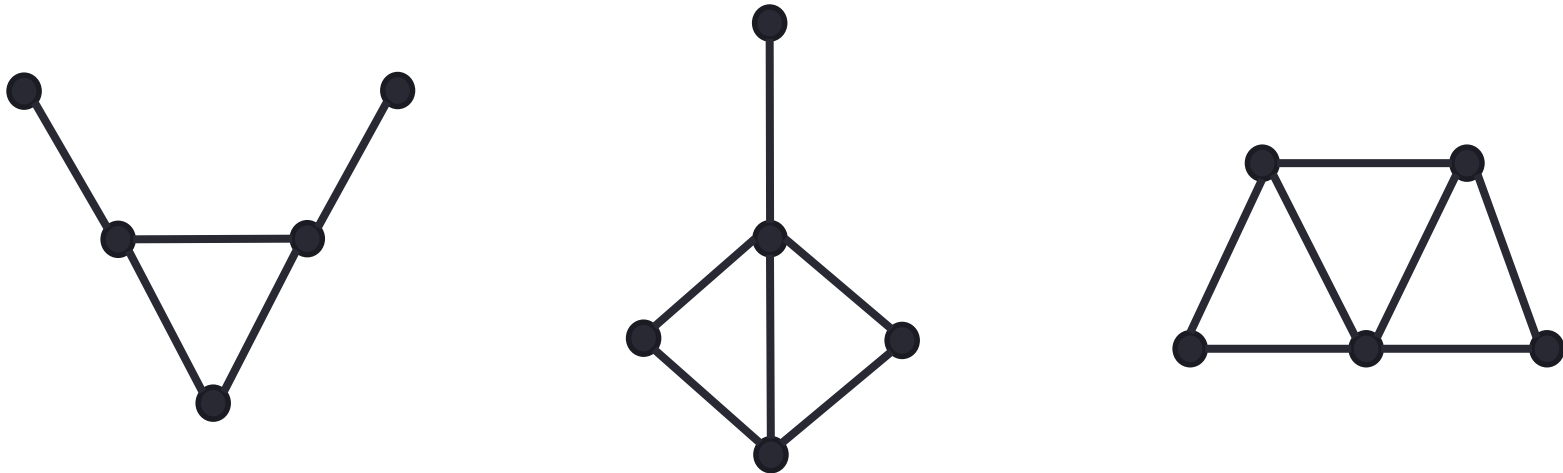
So what do we know about leaf
powers?

Proposition

A graph G is a 2-leaf power if and only if it contains no induced P_3 (a path of length 2).

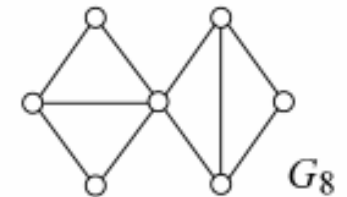
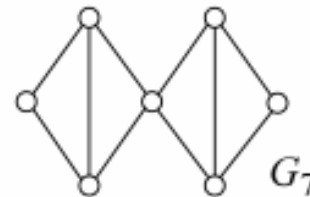
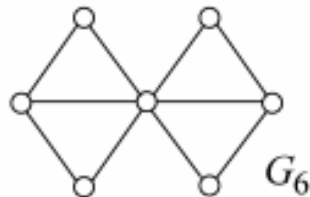
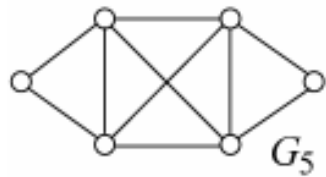
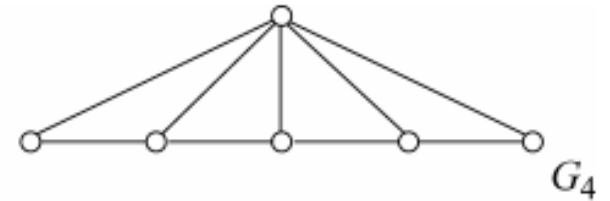
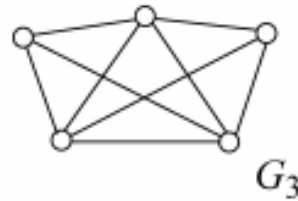
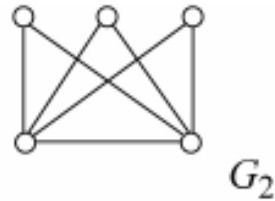
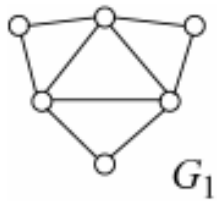
Theorem [Dom, 2006][Rautenbach, 2006]

A graph G is a 3-leaf power if and only if it is **chordal** and does not contain a bull, dart or gem as an induced subgraph.



Theorem [Brandstädt and Hundt, 2008]

A graph G is a (twin-free) 4-leaf power if and only if it is **chordal** and does not contain any of these graphs:



Algorithmic results

Theorem [Chang & al., 2007]

5-leaf powers can be recognized in polynomial time.

Theorem [Eppstein, Havvaei, 2018]

k -leaf powers with bounded degeneracy d can be recognized in time $f(k, d)poly(n)$.

Theorem [Ducoffe, 2018]

6-leaf powers can be recognized in polynomial time.

[65 pages]

Results on leaf powers

- If G is a leaf power, then G is chordal (i.e. every cycle of length at least 4 has a shortcut).

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 - Idea: taking the power of a tree does not create chordless cycles.

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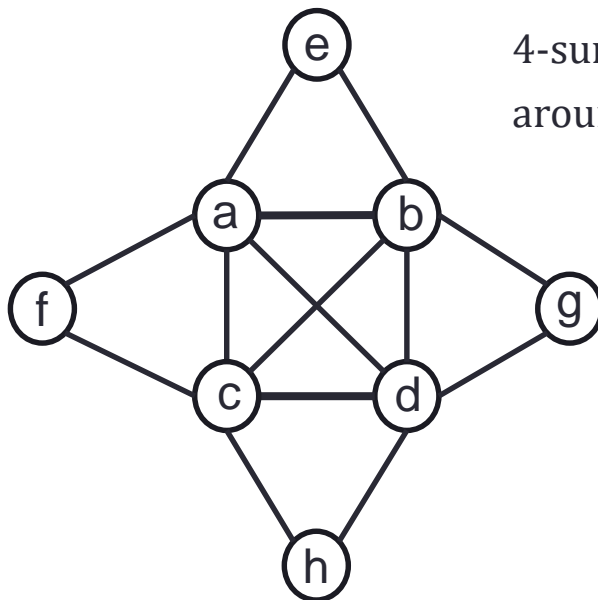
- If G is a leaf power, then G is **chordal** (i.e. every cycle of length at least 4 has a shortcut).
 - Idea: taking the power of a tree does not create chordless cycles.
- If G is a leaf power, then G is **strongly chordal**.

Results on leaf powers

- If G is a leaf power, then G is **strongly chordal**.

A graph is strongly chordal if it is chordal and k -sun-free, for every $k \geq 3$.

A k -sun is the graph obtained by starting with a clique with vertices v_1, \dots, v_k , and adding a common neighbor to each v_i, v_{i+1} pair.



4-sun: start with a K_4 , add 'spikes' around the clique.

Open question (at the time)

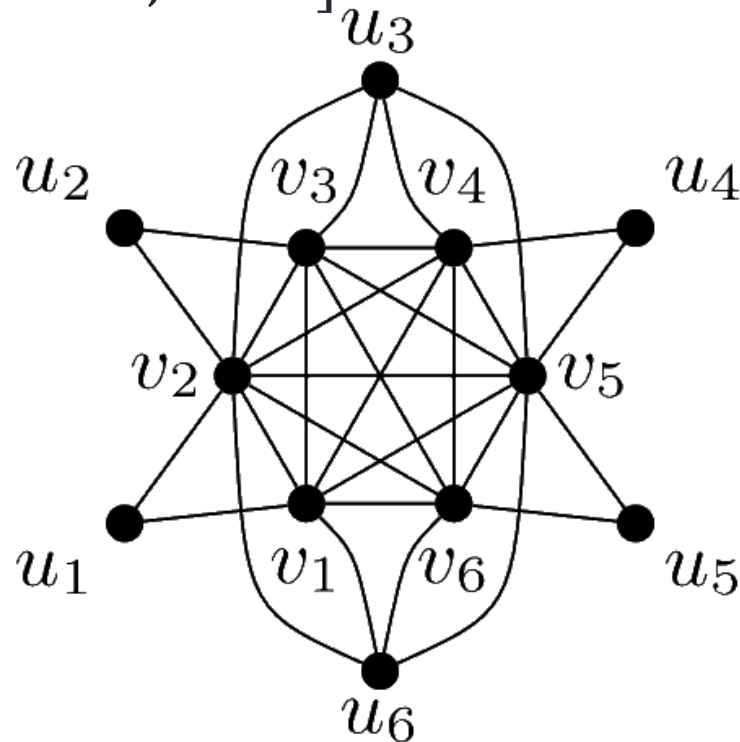
Are leaf powers exactly the strongly chordal graphs?

Open question (at the time)

Are leaf powers exactly the strongly chordal graphs?

NO

[Brandstädt and Hundt., 2010]



Strongly chordal,
not a leaf power

Open question

Is there only one (minimal) strongly chordal graph that is not a leaf power?

Minimal in the sense that removing a vertex makes the graph a leaf power.

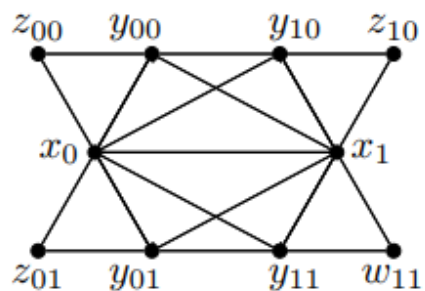
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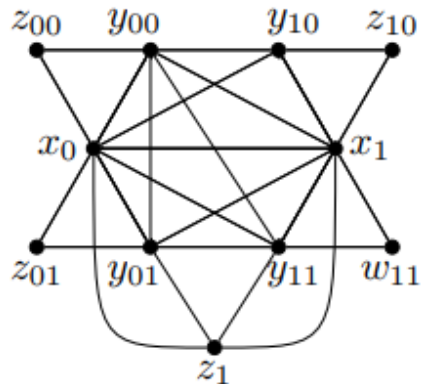
NO

There are at least 7 of them!

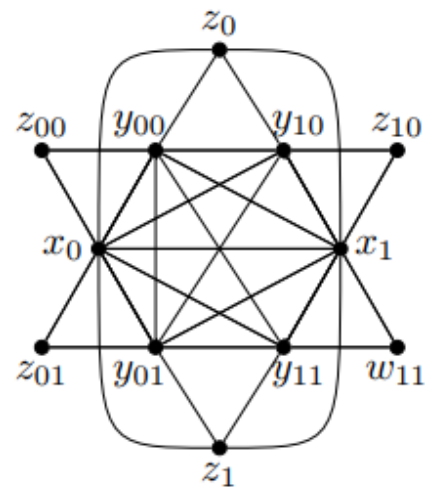
[Nevries & Rosenke, 2015]



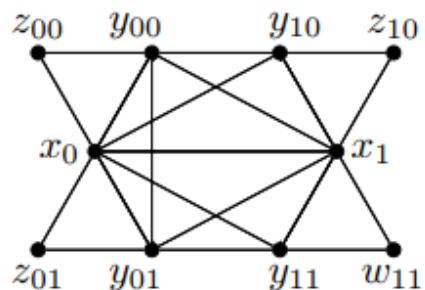
G_1



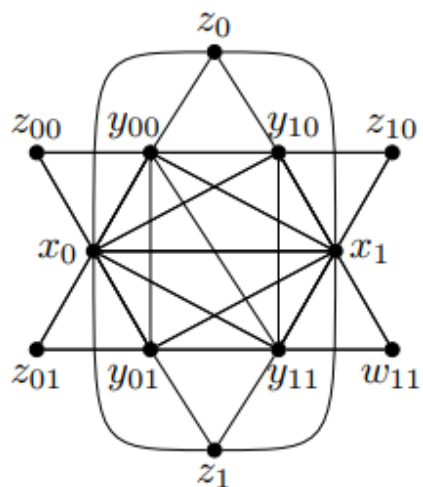
G_4



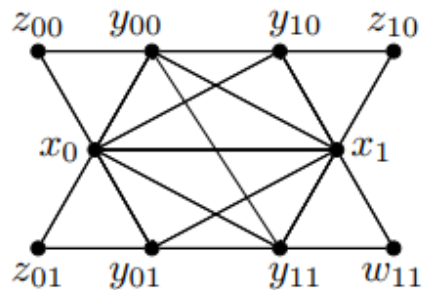
G_6



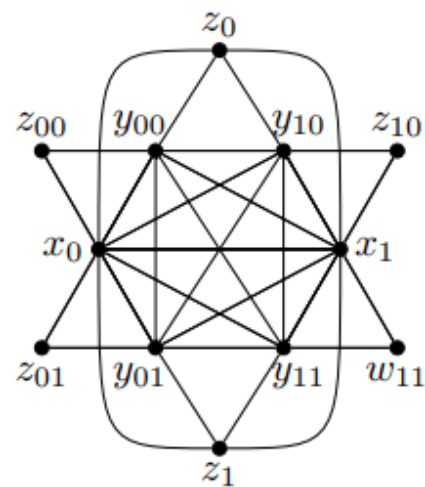
G_2



G_5



G_3



G_7

Conjecture

There are only a finite number of (minimal) strongly chordal graphs that are not leaf powers.

Conjecture

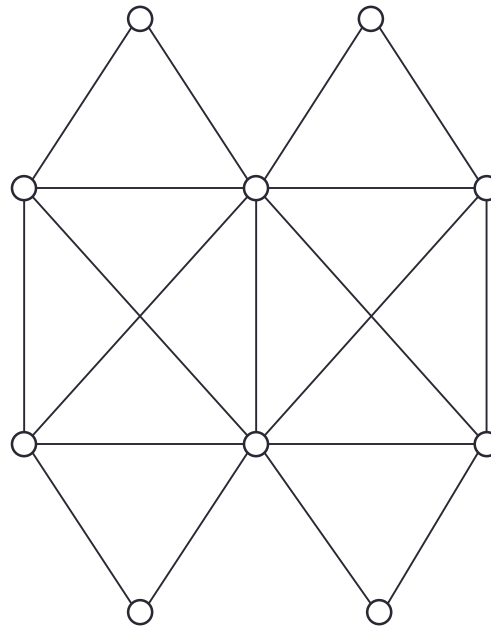
There are only a finite number of (minimal) strongly chordal graphs that are not leaf powers.

NO!

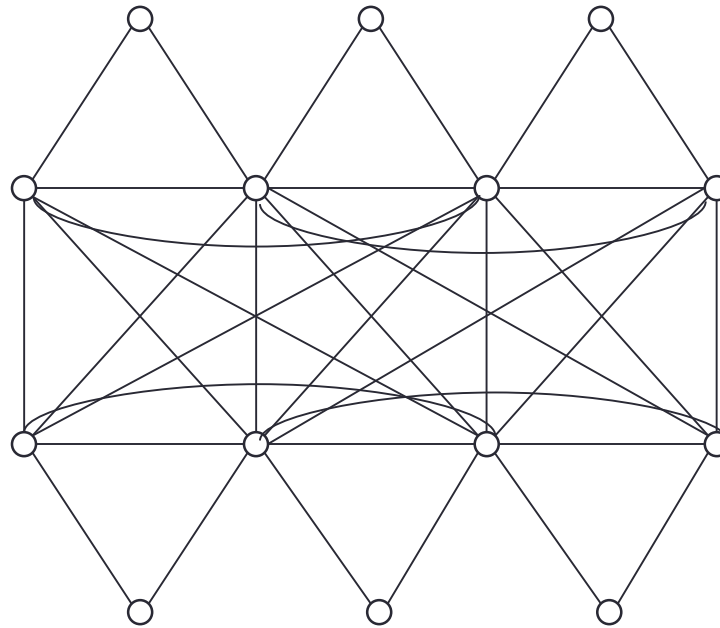
Theorem [L, 2018]

There is an infinity of (minimal) strongly chordal graphs that are not leaf powers.

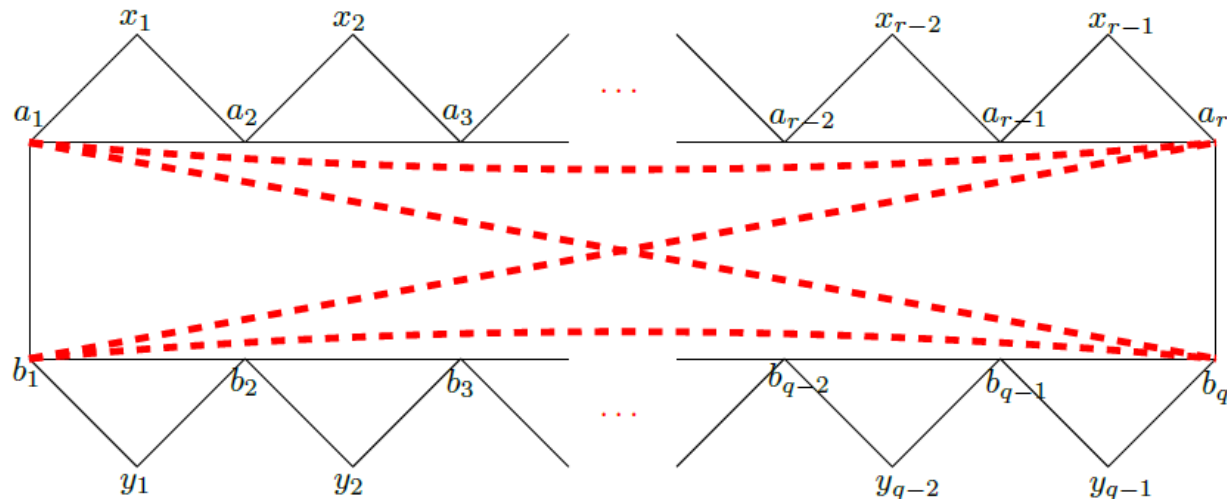
- Here's one of them



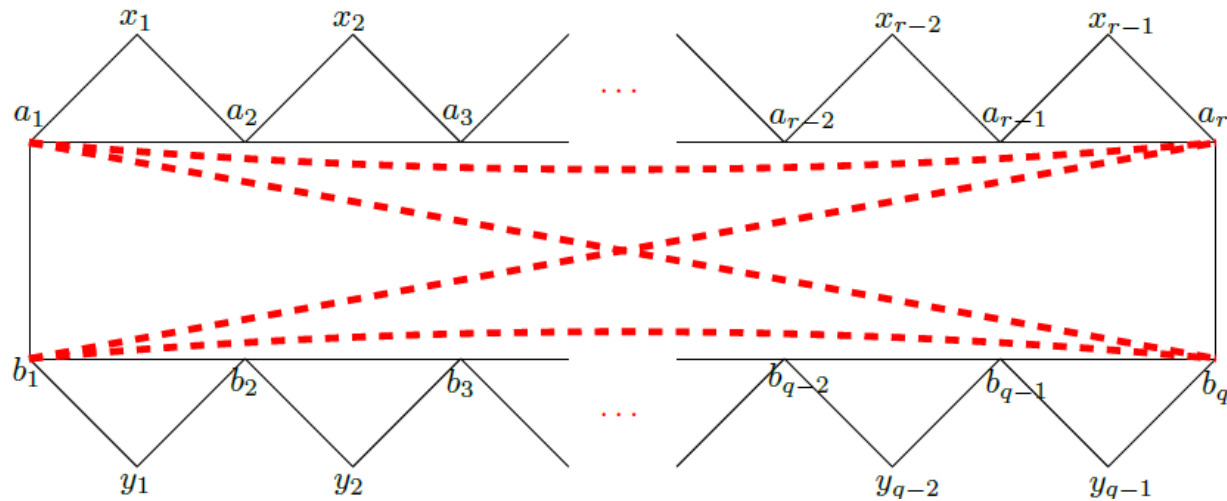
- Here's another



- In general: start with paths of length r and q , add spikes on edges, add all edges between paths vertices except four extremities.



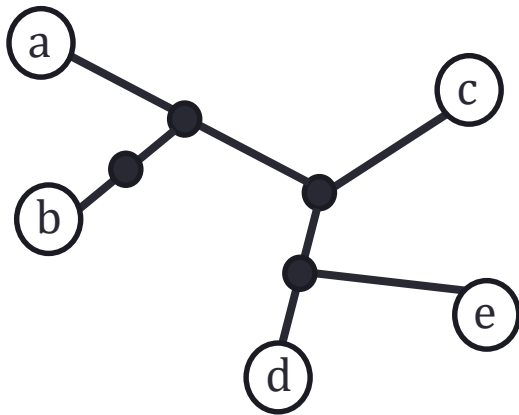
- In general: start with paths of length r and q , add spikes on edges, add all edges between paths vertices except four extremities.



- Not too hard to prove strong chordality.
- But how do we prove non-leaf powerness?

Detecting non-leaf powers from quartets

- Mandatory quartet consistency
 - A new way of identifying non-leaf powers
- **Definition:** a tree T contains a quartet $ab|cd$ if T has an edge that separates $\{a,b\}$ from $\{c,d\}$



Contains:

$ab|cd$

$ab|ce$

$ab|de$

$ac|de$

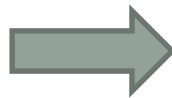
$bc|de$

Assume that G is a leaf power explained by some tree T .

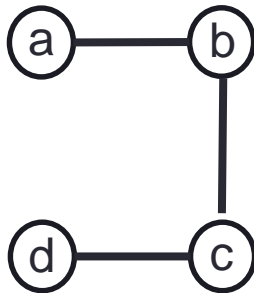
If G has one of these...



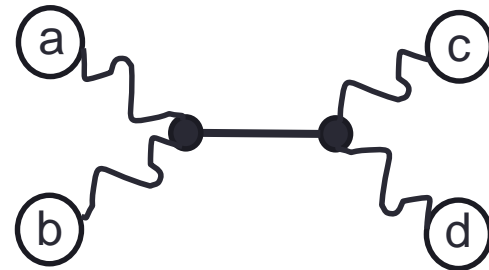
$2K_2$



P_4



..then T must contain the $ab|cd$ quartet.



Detecting non-leaf powers from quartets

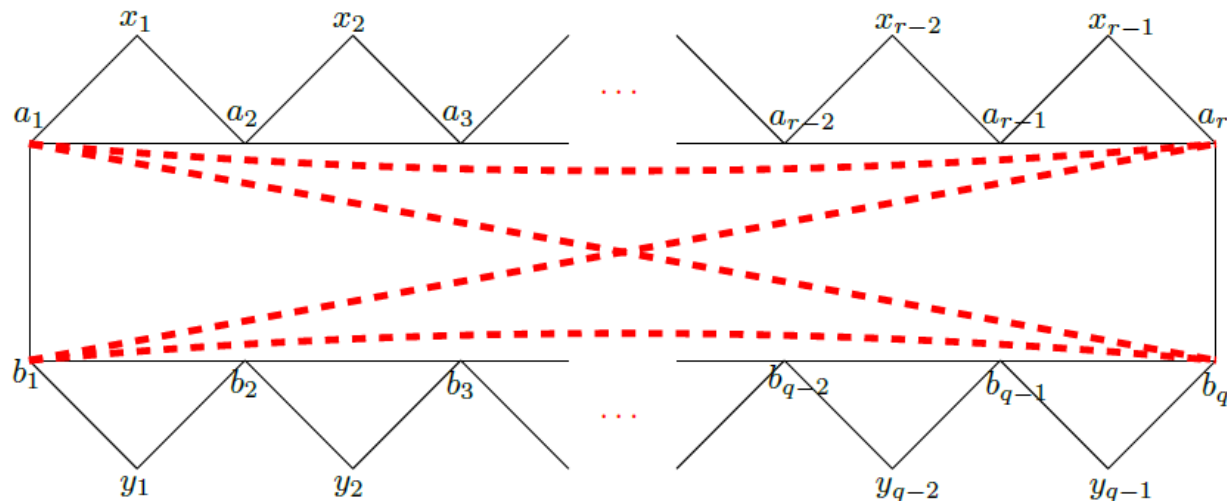
- List every $2K_2$ and P_4 to infer the mandatory quartets Q .
- If no tree can possibly contain every quartet of Q , then G cannot be a leaf power.

Detecting non-leaf powers from quartets

For these graphs, can be shown to require

every $a_i a_{i+1} | b_j b_{j+1}$, and $a_1 b_1 | a_r b_q$

which are known to admit no tree [Shutters et al., 2012]



Some trivia

- NP-hard to decide if a chordal graph contains a copy of some $G_{r,q}$
 - First non-leaf power certificate that we don't know how to check
- Conjecture: can recognize k -leaf powers explained by a $f(k)$ bounded degree tree in $\text{poly}(n)$ time.
 - Implications?
- Other necessary conditions based on alternating cycles (edge to non-edge alternation)

Conclusion

- Every question is still open
 - Complexity of recognizing leaf powers and k -leaf powers?
 - Forbidden subgraphs of leaf powers and k -leaf powers?
- Leaf powers and quartets
 - Question: is G a leaf power iff there is a tree that contains all its mandatory quartets?
 - Answer: no, but only one counter-example is known, the 3-sun. All other non-leaf powers have incompatible quartets.