ALGORITHMIC CHALLENGES IN RECONSTRUCTING COPY-NUMBER EVOLUTION

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Tumours and evolution



- Tumor cells undergo rapid genomic evolutionary changes
 - Rearrangements
 - Amplifications
 - Deletions
 - e.g. : breakage-fusion-bridge cycle















Single cell sequencing

- Assessment of tumor heterogeneity
- Cells within a tumor undergo somatic mutations
 => Different rearrangements, duplications, deletions
 => Difficult to obtain complete genome of every cell
- Copy-numbers are easier to obtain.
 - High coverage = expensive \$\$
 - Accuracy at low coverage now possible (Direct libray preparation, 10X Genomics, ...)

Copy-Number Profile (CNP)



10 copies of segment 1
5 copies of segment 2
0 copies of segment 3
4 copies of segment 4
10 copies of segment 5

Copy-Number Profile (CNP)



(10, 5, 0, 4, 10)



Our Solution Features

- Profiles 100s to 1000s of cells
- Accurately calls single cell CNV events at 2 Mb resolution
- · Detects CNV events down to 100s of Kb on clusters of cells
- Demonstrated with cell lines, primary cells, fresh and frozen tissue
- Push-button analysis and visualization software



View the Workflow

Sequence cells from same tumor (single cell sequencing) Infer copy-number for each segment of interest (for each allele) Phasing: assign copy-number to allele, get chromosome Copy-Number Profile (CNP) Compare CNPs of each pair of cells =>distance matrix Reconstruct phylogeny





Major: (4, 4, 12, 0, 2) Minor: (4, 3, 10, 0, 1)

0	16	47	72	77	79
16	0	37	57	65	66
47	37	0	40	30	35
72	57	40	0	31	23
77	65	30	31	0	10
79	66	35	23	10	0











If we know the Copy-Number Profile (CNP) of each segment of interest in several tumor cells, what can we say about their evolution?

In this talk

• Comparing integer vectors.

Comparing genomes an CNPs.

• Comparing integer vectors (with rearrangements).

In this talk

- Comparing integer vectors.
 - We have no idea how!
- Comparing genomes an CNPs.
 - We have no idea how!
- Comparing integer vectors (with rearrangements).
 - We have no idea how!

Comparing Copy-Number Profiles

Why compute distances

- Classic approach 1
 - compute dist(u, v) for each pair u, v
 - get a distance matrix
 - use phylogenetic distance method (e.g. NJ)

	X	У	Z
X	5	8	6
у		5	7
z			2

- Classic approach 2
 - infer ancestral CNP states
 - minimize sum of branch distances



The story so far

- 1. Let's use the Euclidean distance (2011)
- 2. Let's model segmental events on integer vectors (2014)
 - Even if minimizing events takes exponential time...
 - No actually it takes polynomial time (2017)
- 3. Let's weight events by their amplitude (2019)
- 4. Let's weight events by their length/location (upcoming...)



Letter Published: 13 March 2011

Tumour evolution inferred by single-cell sequencing

Nicholas Navin, Jude Kendall, Jennifer Troge, Peter Andrews, Linda Rodgers, Jeanne McIndoo, Kerry Cook, Asya Stepansky, Dan Levy, Diane Esposito, Lakshmi Muthuswamy, Alex Krasnitz, W. Richard McCombie, James Hicks & Michael Wigler ™

 Sequenced 100 cells from a tumor, reconstructed NJ phylogeny from CNP data.



- In Navin et al. [Nature11] : Euclidean distance
- $dist(\boldsymbol{u}, \boldsymbol{v}) = \sqrt{\sum (u_i vi)^2}$



dist = $\sqrt{9 + 16} = 5$

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- $dist(\boldsymbol{u}, \boldsymbol{v}) = \sqrt{\sum (u_i vi)^2}$



Implicit assumption: positions are independent.

In Schwarz et al. [PlosCB14]: MEDICC model

- Positions should NOT be independent!
- Events can affect segments of genomes

CNP-2-CNP problem – MEDICC model

- <u>**Given</u>**: two CNPs u and v (integer vectors)</u>
- <u>Move</u>: alter an interval of u by +1/-1 (a 0 stays a 0).
- **<u>Find</u>**: min # of moves to turn u into v

(1,1,1,1,1,1,1,1)

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$$(1,1,0,0,1,0,1)$$

$$(1,1,0,0,1,0,1)$$

$$(1,2,0,0,2,0,2)$$

CNP-2-CNP problem – MEDICC model

- In Schwarz & al [PlosCB14]
 - Compute d(u, v) in time $\Omega(3^N)$, N = max copy-number



ZZS algorithm

- In Zeira, Zehavi, Shamir [JCB17]:
 - Better algorithms to compute min # of +1/-1 moves
 - Simple DP pseudo-polynomial time algorithm O(nN²)
 - More involved O(n) time algorithm.

ZZS algorithm

- In Zeira, Zehavi, Shamir [JCB17]:
 - Better algorithms to compute min # of +1/-1 moves
 - Simple DP pseudo-polynomial time algorithm O(nN²)
 - More involved O(n) time algorithm.
- General idea:
 - Show that some optimal solution does all deletions before amplifications.
 - Dynamic programming, optimal for every prefix from left to right.
 - M[i, d] =optimal for *i* if *i*-th value is *d*.

 $\mathbf{M}[i,d] \leftarrow \min_{0 \le d' \le N} \{ \mathbf{M}[\operatorname{prev}(i),d'] + \max\{d-d',0\} + \max\{a(i,d) - a(\operatorname{prev}(i),d'),0\}$

 $+ \max\{Q_i - \max\{d, d'\}, 0\}\}.$

At Recomb-CG 2019

Same problem, but not restricted to +1/-1.

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Same problem, but not restricted to +1/-1.

A single event could double copy numbers (e.g. WGD).

CNP-2-CNP problem – extended MEDICC

- **<u>Given</u>**: two CNPs u and v (integer vectors), cost function f
- <u>Move</u>: alter a contiguous interval of u by any amount.
- **<u>Find</u>**: min # of moves to turn u into v

CNP-2-CNP problem – extended MEDICC

• Difference vector
$$w = u - v$$

Intuition: "squish" values of w to 0.



Theorem

In the extended MEDICC model, the CNP-2-CNP problem is **strongly** NP-hard.

(strongly => hard even if the numbers are polynomial in n)

Positive results

Theorem

If the CNP's have no 0-positions, there is a linear time factor 2 approximation algorithm for the extended CNP-2-CNP problem.

The algorithm

Return the number of flat intervals in the difference vector.

Flat interval = contiguous positions in which difference vector has same value.

Below: 5 flat intervals

<u>Lemma</u>

One moves reduces number of flat intervals by at most 2 => $dist_{any}(u, v)$ is at least ½ the number of flat intervals.



Flat interval = contiguous positions in which difference vector has same value.

Below: 5 flat intervals

<u>Lemma</u>

One moves reduces number of flat intervals by at most 2 => $dist_{any}(u, v)$ is at least ½ the number of flat intervals.

Trivial 2-approx: remove each flat interval one by one.


Experiments

 If we simulate amplifications and deletions on genomes (and not CNPs), can we reconstruct phylogenies?





More leaves = easier to predict



Extended-extended MEDICC model

- Some events are more likely to affect certain regions of the genome.
 - e.g. arm duplications => ends of CNP vector more susceptible to amplification
- Extended model : each interval [i ... j] has its own weight.
- Weights can be inferred from cancer patient data.
- (not published yet, and not my work)

Comparing Genomes with Copy-Number Profiles



Problems with segmental events on CNPs

- Assumes that order of segments remains fixed.
- Rearrangements change the order.
 - Some, drastically.

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- Rearrangements change the order.
 - Some, drastically.

Introduce actual rearrangements into the model.

- In [Zhu & al, ACM-BCB 2018]:
- Given the CNP of of a single cell *C*, infer the rearrangements that occurred from a healthy genome to *C*.



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- Given the CNP of of a single cell *C*, infer the rearrangements that occurred from a healthy genome to *C*.
- Allowed: segmental duplications & deletions.



- **<u>Given</u>**: string *S* and copy-number vector *C*
- Move: segmental duplications and deletions.
- **<u>Find</u>**: min # of moves to turn *S* into any *T* whose CNP is *C*





(2, 4, 1, 3)

Need	
2 x a	
4 x b	
1 x c	
3 x d	



a<u>bcd</u>

a<u>bcdbcd</u>

Segmental duplication (this one is tandem)

(2,	4,	1,	3)
	,	,	

Need	
2 x a	
4 x b	
1 x c	
3 x d	



ab<mark>e</mark>dbcd

abdbcd

(2, 4, 1, 3)

Deletion (this one is not segmental)

Need	
2 x a	
4 x b	
1 x c	
3 x d	



abcdbcd

<u>abdb</u>cd

abdbabdbcd

(2, 4, 1, 3)

Need	
2 x a	
4 x b	
1 x c	
3 x d	



abcdbcd

abdbcd

abdbabdbcd

(2, 4, 1, 3)

Need 2 x a 4 x b 1 x c 3 x d





abcdbcefabbc







abcdbcefabbc





abcdbcefabbc





- The problem is NP-hard [Zhu & al., 2018]
- Reduction from set-cover: design S and CNP C so that
 - each elements = 1 character
 - choosing a set = deleting elements
 - must delete one occurrence of each element

In [Lafond, Zhu & Zou, CPM, submitted]

<u>Theorem</u>

The Genome-to-CNP problem (probably) does not admit a constant factor approximation and (probably) is not FPT.





Open problem

Find any practical approach!

Open problem

If initial string *S* is exemplar, is Genome-to-CNP in P?

Exemplar = no characer occurs more than once.

Could be useful: we may model each chromosome of the healthy genome as exemplar.

Comparing Integer Vectors (with rearrangements)



Interval events may give rise to impossible scenarios.

$$(1,1,\underline{1,1},1,1,1)$$

$$(1,1,0,0,1,\underline{1},1)$$

$$(1,1,0,0,1,0,1)$$

$$(1,1,0,0,1,0,1)$$

$$+1$$

$$(1,2,0,0,2,0,2)$$

- Interval events may give rise to impossible scenarios.
- Compare CNPs, but require the existence of actual genomes + rearrangements.

$$(1,1,\underline{1,1},1,1,1)$$

$$(1,1,0,0,1,\underline{1},1)$$

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$$(1,2,0,0,2,0,2)$$

- <u>Given</u>: two CNPs u and v
- <u>Move</u>: segmental duplications and deletions (on genomes).

• <u>Find</u>:

- a genome G_1 whose CNP is \boldsymbol{u} ;
- a genome G_2 whose CNP is v;
- such that # of dups/deletions from G_1 to G_2 is minimum.

(1, 2, 2, 2)

Go from any genome with 1 x a 2 x b 2 x c 2 x d to any genome with 3 x a 0 x b 3 x c 6 x d

(3, 0, 3, 6)

(1, 2, 2, 2)

addccbb

(3, 0, 3, 6)

Go from any genome with 1 x a 2 x b 2 x c 2 x d to any genome with 3 x a 0 x b 3 x c 6 x d

Go from any (1, 2, 2, 2)genome with addccbb 1 x a 2 x b (1, 0, 2, 2)2 x c addcc 2 x d to any genome with 3 x a 0 x b (3, 0, 3, 6)3 x c 6 x d

Go from **any** (1, 2, 2, 2)genome with 1 x a addccbb 2 x b(1, 0, 2, 2)2 x c <u>addc</u>c 2 x d addcaddcc (2, 0, 3, 4) to any genome with 3 x a 0 x b (3, 0, 3, 6)3 x c 6 x d

(1, 2, 2, 2)

addcc<mark>bb</mark>

<u>addc</u>c (1, 0, 2, 2)

 \underline{add} c a d d c c (2, 0, 3, 4)

<u>add</u>addcaddcc

(3, 0, 3, 6)

Go from **any** genome with 1 x a 2 x b2 x c 2 x d to any genome with Зха 0 x b 3 x c 6 x d
Consistent CNP-2-CNP problem

Open problem

Any question you can think of about this problem!

- Very interesting theoretical problem.
- In practice...
 - Some optimal solution always has 0 or 1 deletion (we think)
 - Gives rise to ridiculous genomes
 - e.g. aaaaaabbbbbbbcccccdddd
- More useful formulation: global inference of genomes on a phylogeny

Phylogenetic CNP problem

- <u>Given</u>: phylogeny T with CNPs at leaves, human genome at root
- **Find**: a genome assignment at each node of T such that:
 - each genome at a leaf has correct CNP;
 - sum of rearrangements at branches is minimum.



(10, 5, 0, 4, 10) (8, 6, 2, 0, 8) (8, 6, 2, 0, 8)

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- <u>Given</u>: phylogeny T with CNPs at leaves, human genome at root
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- Copy-number profiles carry useful information on tumor heterogeneity.
- Easier to obtain than whole genomes.

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- Copy-number profiles carry useful information on tumor heterogeneity.
- Easier to obtain than whole genomes.
- We are not exploiting this information at its full potential!

Conclusion

• What this line of research needs:

- Better models
- Better problem formulations
- Better algorithms
- Better access to real data!!!