Algorithms for the validation and correction of orthology relations

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Introduction

Gene trees, species trees Duplication, speciation Orthologs, paralogs, why?

Validation and correction of orthology relations Cograph (P₄-free) characterization of valid relations Modeling uncertain relations

Similarity graphs vs orthology graphs

- Why they are not the same
- How to deal with similarity graphs



Introduction

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Similarity graphs vs orthology graphs



Take some gene, say my favorite RPGR :Retinitis pigmentosa GTPase regulatorParticipates in eye coloring.

What is the **history** of RPGR ?

Almost all vertebrates have a copy of this gene. Some have more than one. Some don't have it.

What happened exactly?

A gene can be :

- Transmitted to descending species by speciation
- Duplicated
- Lost



































Orthologs and paralogs

Two genes are*:

Orthologs if their lowest common ancestor underwent **speciation**

Paralogs if their lowest common ancestor underwent **duplication**

*w.r.t. a given gene tree







Why bother?

Orthology/paralogy relations are related to gene functionality.

Some gene functional annotation databases assume that orthologs share the same functionality. (e.g. COG, eggNOG databases)

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Orthologs conjecture: orthologous genes tend to be similar in function, whereas paralogous genes tend to differ.

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Quest For Orthologs consortium: "a joint effort to benchmark, improve and standardize orthology predictions through collaboration, the use of shared reference datasets, and evaluation of emerging new methods".

Clustering genes into groups of orthologs:

- If g1 and g2 and "**similar enough**" in terms of sequence, we say that g1 and g2 are putative orthologs.
- Make a graph G of putative orthologs.
- Partition G into clusters, i.e. highly connected components Otherwise, too many false positives occur
- OrthoMCL, InParanoid, proteinortho, ...



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Clustering genes into groups of orthologs:

- If g1 and g2 and "similar enough" in terms of sequence, we say that g1 and g2 are putative orthologs
- "Similar enough" usually means that, if g1 and g2 are from species s1 and s2, they for a Bidirectional Best Hit (BBH):
 - g1's best match in s2 is g2
 - g2's best match in s1 is g1



These methods are very often **incomplete** - have false positives or false negatives (according to our definitions).

In (Lafond & El-Mabrouk, 2014), we found that >70% of inferred sets of relations were **unsatisfiable** – corresponded to no possible gene tree.



Orthology/paralogy relation graph R



Orthologs = (a,b) (a, c) (c, d)Paralogs = (a, d) (b, c) (b, d)

OrthologsParalogs

Orthology/paralogy graph Orthologs = (a,b) (a, c) (c, d) Paralogs = (a, d) (b, c) (b, d)



What we want to do

Given a set of orthologs / paralogs in form of a relation graph R:

Verify that they "make sense"

Satisfiable: can some gene tree display the relations? Consistent: does it agree with our species tree?

 If they don't make sense, correct them in some minimal way Everything is NP-Complete






Problem :

Given a relation graph R, is R satisfiable?

Does there exist a gene tree G that displays the relations of R ?



Usages of verifying satisfiability

- 1. Orthology graph benchmarking
- 2. Gene tree reconstruction
- 3. Species tree reconstruction



So, how do we verify whether there is a gene tree displaying these relations?

And if so, can we construct the tree?



Theorem (Hernandez-Rosales & al., 2012):

A relation graph R is satisfiable if and only if R_{BLACK} is P_4 -free (has no induced path on 4 vertices).

(P₄-free graphs are sometimes known as cographs)



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Is there a gene tree for R?



Let's say it exists...what is the first split then ?















Monochromatic edge-cut => a split exists



and so on ...

Theorem (informal) (Corneil, Perl & Stewart, 1985)

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 P_4 -freeness is easy to check in polynomial time. O(n⁴) in the obvious way, O(n) in more clever ways.



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We looked at 265 inferred families from **ProteinOrtho**, under 5 parameter sets {-2, -1, 0, +1, +2}.

Stricter => Less orthologies



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Unknown/undecided relations

We might lack confidence in some given relations e.g. genes having a borderline BLAST similarity value



Problem :

Given a relation graph R with **unknown edges**, can they be **chosen** to make R:

- satisfiable?
- S-Consistent?
- self-consistent?



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Can we get some robust relationships out of these ?

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Can we get some robust relationships out of these ?



Keep the common orthologies and paralogies.

The rest is unknown.





Gene relation correction

Make R satisfiable by changing a minimum number of relations.

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NP-Complete (El-Mallah & Colbourn, 1988)



- Many other variants, all difficult:
 - Remove as few genes to have a P4-free graph => can't even approximate
 - Incorporate information from species tree => still NP-complete
 - Add weights on the orthology/paralogy relations => can't approximate
 (Dondi, Lafond, El-Mabrouk, 2014-2016)

ILP formulation (has difficulty handing > 10 genes)
FPT algorithms (also slow)
MinCut heuristic (no performance guarantees)

Dealing with similarity-based methods

Orthology/paralogy relation graph R

Sequences and stuff



R a c c d

Orthologs = (a,b) (a, c) (c, d)Paralogs = (a, d) (b, c) (b, d)

> Orthologs Paralogs

Orthology/paralogy relation graph R

Sequences and stuff



Traditional inference method

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Orthology/paralogy relation graph R









Similarity graph



Similarity graph



Similarity graph

The (a, b2) orthology is indistinguishable from paralogy from the point of view of similarity.



BAD for:

- 1) Benchmarking: the graph passes the test of being P4free, and yet does not depict relations correctly
- 2) Gene tree reconstruction: interpreting as relations yields the wrong gene tree.

- 1) Give up on these missing orthologs.
- 2) Devise methods that really infer relation graphs.
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 - Recognizing leaf-powers is a longstanding open problem (not known to be in P nor NP-complete)
 - Too complicated, let's start with a restricted model

Orthologs conjecture: orthologous genes tend to be similar in function, whereas paralogous genes tend to differ.

- 1) In the absence of gene duplication, no significant dissimilarity should be observed.
- 2) In the event of gene duplication, one copy remains intact whereas the other evolves at an accelerated rate.

(as in the motivation for the orthologs conjecture)



Direct consequences of the axioms of the DAD model:

- Two genes will appear as "non-similar" if and only if a divergent duplication edge separates them.
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- Clustering algorithms can be applied to find the "similarity cliques", which we assume represent orthology subtrees.
- The cliques do not represent all orthologies: some (and perhaps many) may be missing, e.g. (b, f), (b, g), (c, f), ...





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- How can we find missing relations?
 - (WIP)



Conclusion

- Orthology/paralogy graphs are exactly the P₄-free graphs
- In practice, we only have a similarity graph
 - Not the same
- Can we "turn" a similarity graph into an orthology/paralogy graph?
 - What are the limits of similarity for orthology inference?
- Future works: design algorithms to infer missing orthologs from a similarity graph, and test them on real/simulated datasets.

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NP-Complete (Lafond & El-Mabrouk, 2014)

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NP-Hard to approximate within a n^{1-ε} factor. (Lafond, Dondi, & El-Mabrouk, 2016)

Weighted gene relation correction

To make things easier:

Give each edge a **weight**, representing some degree of confidence over the inferred orthology/paralogy.

This weight represents the cost for changing the edge's color.



Weighted gene relation correction

Something we can handle:

If edges all have weights of 0 or 1

0 = don't care, 1 = don't touch

We can tell in polynomial time if there is an edge editing of weight 0.



Weighted gene relation correction

If weights are arbitrary, NP-Hardness follows from the unweighted version (for both satisfiability and consistency).

Worse than that, there is **no constant factor approximation** assuming the *unique games conjecture*.



Fixed parameter tractability

k = number of edges that can be edited

For **satisfiability**, the unweighted edge-editing problem admits a vertex kernel of size O(k³) (Guillemot, Paul, Perez, 2010)

There is an obvious FPT algorithm:

each P_4 must be killed. There are 6 edge modifications that accomplish this. Branch into each possibility. O(6^kn)

• can be extended to S-consistency

Was improved to O(4.612^k + |V|^{4.5}) (Lui, Wang, Guo & Chen, 2012)

• good for S-consistency? No idea.
Y. Liu et al. / Theoretical Computer Science 461 (2012) 45-54



Min-cut approximation for satisfiability

Recall:

Theorem (again):

A relation graph R is satisfiable if and only if for each subgraph R', one of R'_{BLACK} or R'_{BLUE} is disconnected.

In particular, R_{BLACK} or its complement R_{BLUE} must be disconnected.

So we'll disconnect it then.

Min-cut approximation for satisfiability

In particular, R_{BLACK} or its complement R_{BLUE} must be disconnected.

Find a min-cut on R_{BLACK} Find a min-cut on R_{BLUE} Take the best of the two and apply. Repeat on the resulting components.

(min-cut = minimum weight edge-set that disconnect R, can be found in time $O(n^3)$)

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Gives a solution that is at most n times worse than optimal. (not great, but shows that approximability is bounded)

Theoretical and practical problems

Theoretical problems

Unweighted case: can we approximate satisfiability? Consistency?

Weighted case: gap in approximability results. Is there better than a n-factor approximation? Somewhere in-between constant and n.

FPT : elements of unweighted satisfiability correction (aka cographediting) are known. Not much about the rest.

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(but earlier I said we could!)

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Similarity graph =/= orthology/paralogy graph



We don't even know **how to test** our correction methods.

Gold standard datasets are extremely rare, if nonexistent. Most software are interested into forming clusters of orthologs. How do we compare with others?

Faster approximations and heuristics are still needed.

The Min-Cut algorithm takes time $O(n^3)$, and our implementation is too slow for, say, 1000 genes.

How to handle **other events**?

How can we distinguish species tree disagreement with HGT or ILS? Beyond graph theory, what is their practical impact in the ortholgoy/paralogy inference process?